

Note Every subgraph of a planar graph is planar

Note Every graph which has a non-planar subgraph is also non-planar

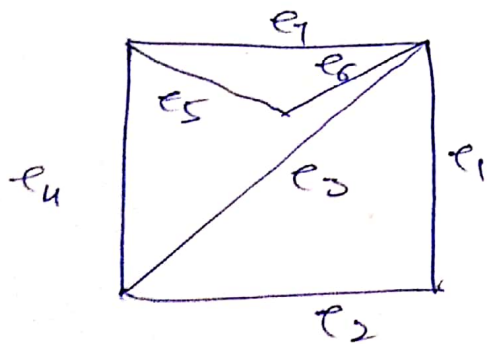
- A disconnected graph is planar iff each of its components is planar

Note A separable graph (v.c. is one) is planar iff each of its blocks is planar.

- Addition or removal of self-loops in a graph does not affect planarity

- // edges do not , , .

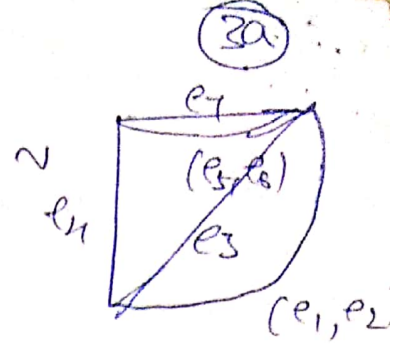
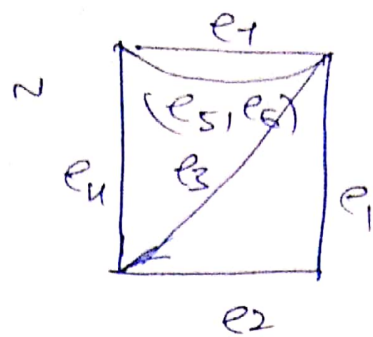
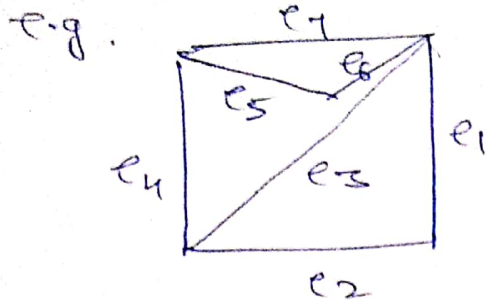
Defn → In a graph, two edges are said to be in series if they have exactly one vertex in common, and if this vertex is of degree 2. e.g.



$e_5 \& e_6$ in series

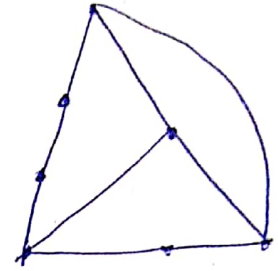
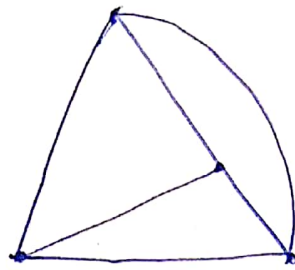
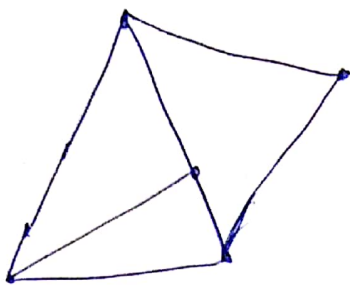
$e_1 \& e_2$ in series

Homeomorphic Graphs → $G_1 \& G_2$ hom. if one graph can be obtained from the other by the creation of edges in series i.e. by insertion of vertices of degree two or by the merger of edges in series.



homeomorphic gphs.

e.g.



\rightarrow reduced edges (or vertex of deg 2)

\rightarrow added edges (or vertex of deg 2)

Note \rightarrow A gph is planar iff every graph that is homeomorphic to it is planar

Note A necc. & suff. cond'n for a gph G to be planar is that it does not contain either of Kuratowski's two gphs or any gph homeo to either of them.

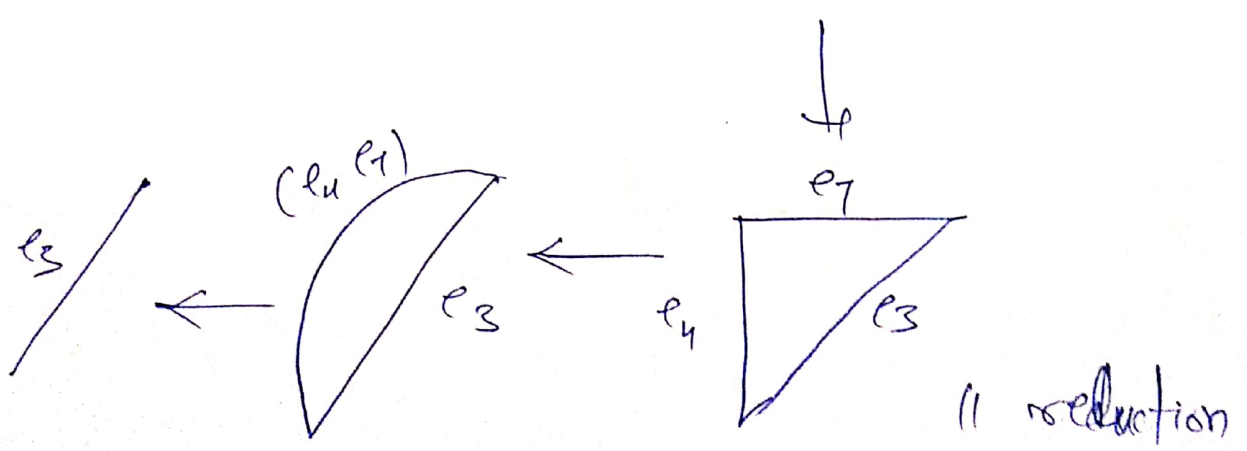
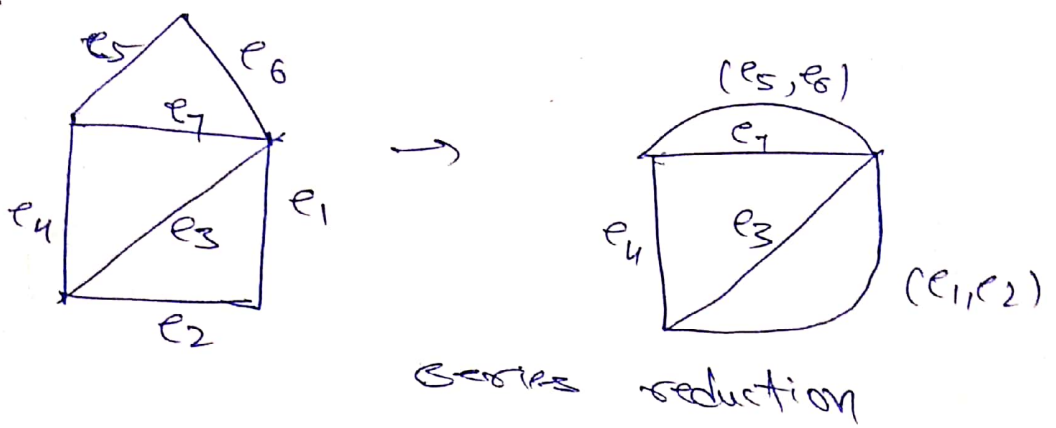
Note A non-planar gph. may have a subgph. homeo to K_5 graph.

Detection of Planarity, How to find if G is planar

1. Remove self-loops
2. " // edges
3. eliminate all edges in a series
repeat 2 & 3 we get a graph H →
This graph H is -

1. A single edge or
2. A complete graph of 4 vertices or
3. A nonseparable, simple gph. with $n \geq 5$ and $e \geq 7$

e.g.



∴ we need to investigate only simple, connected non-sep. gph of at least 5 vertices and with every vertex of deg three.

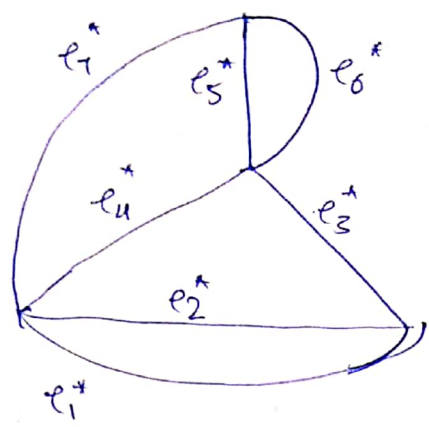
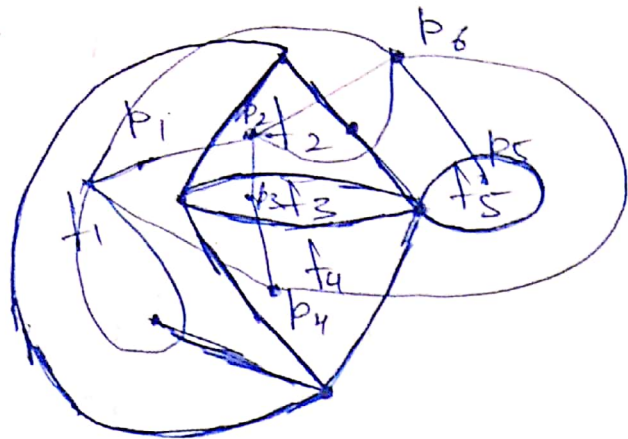
Geometric Dual

- plane representation of a graph (49)
- place one pt. in each region.
- draw edges b/w these pts.
- if two regions are adjacent, draw a line that intersects the common edge
- more than one edge in common then draw one line b/w pts. for each of the common edges.
- for an edge lying entirely in one region, draw a self-loop.

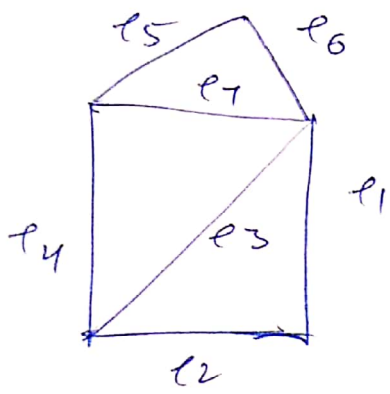
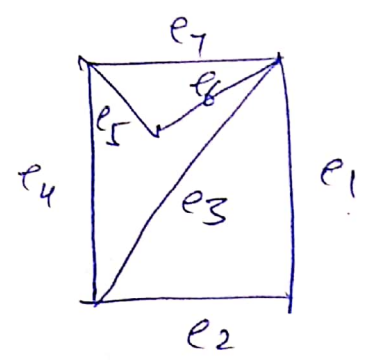
We obtain a graph G^* → geometric dual or Dual of G

- An edge forming a self-loop in G yields a pendant edge in G^*
- A pendant edge in G yields a self-loop in G^*
- edges that are in series in G produce // edges in G^*
- // edges in G produce edges in series in G^*
- G^* is also planar
- G is dual of G^* & vice-versa
- n, e, f, r, u of G
then $n^* = f, e^* = e, f^* = n$
 $r^* = u, u^* = r$

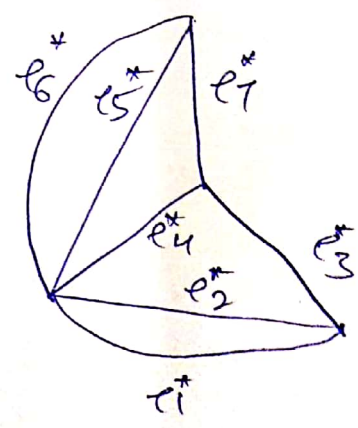
e.g.



Dual
←



Dual →



Self-Dual → If a planar graph is isomorphic to its own dual i.e. $G \cong G^*$
 Complete graph with n vertices is self-dual

Note → $G \rightarrow$ planar graph, $G^* \rightarrow$ Dual
 Then a set of edges in G forms a circuit iff the corresponding set of edges in G^* forms a cut-set in G^* .